

## A Generalized Framework for Auto-tuning Stencil Computations

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# The Challenge: Productive Implementation of an Auto-tuner



- Take one kernel/application
  - Perform some analysis of it
  - Research the literature for appropriate optimizations
  - Implement a couple of them by hand optimizing for one target machine.
  - Iterate a couple of times.
- Result:

improve performance for **one** kernel on **one** computer.



## **Conventional Auto-tuning**

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- Automate the code generation and tuning process.
  - Perform some analysis of the kernel
  - Research the literature for appropriate optimizations
  - implement a code generator and search benchmark
  - explore optimization space
  - report best implementation/parameters
- Result:

significantly improve performance for one kernel on any computer.

*i.e. provides performance portability* 

- Downside:
  - autotuner creation time is substantial
  - must reinvent the wheel for every kernel



## Generalized Frameworks for Auto-tuning

- Integrate some of the code transformation features of a compiler with the domain-specific optimization knowledge of an auto-tuner
  - parse high-level source
  - apply transformations allowed by the domain, but not necessarily safe based on language semantics alone
  - generate code + auto-tuning benchmark
  - explore optimization space
  - report best implementation/parameters
- Result:

significantly improve performance for any kernel on any computer for a domain or motif.

*i.e.* performance portability without sacrificing productivity



## Outline

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- 1. Stencils
- 2. Machines
- 3. Framework
- 4. Results
- 5. Conclusions



# **Benchmark Stencils**

- Laplacian
- Divergence
- Gradient
- Bilateral Filtering



- Nearest neighbor computations on structured grids (1D...ND array)
- stencils from PDEs are often a weighted linear combination of neighboring values
- cases where weights vary in space/time
- stencil can also result in a table lookup
- stencils can be nonlinear operators



 caveat: We only examine implementations like Jacobi's Method (*i.e. separate read and write arrays*)



- ✤ 7-point stencil on scalar grid, produces a scalar grid
- Substantial reuse (+high working set size)
- Memory-intensive kernel
- Elimination of capacity misses may improve performance by 66%





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- 6-point stencil on a vector grid, produces a scalar grid
- Low reuse per component.
- Only z-component demands a large working set
- Memory-intensive kernel
- Elimination of capacity misses may improve performance by 40%



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- 6-point stencil on a scalar grid, produces a vector grid
- High reuse (like laplacian)
- High working set size
- three write streams (+ write allocation streams) = 7 total streams
- Memory-intensive kernel
- Elimination of capacity misses may improve performance by 30%





- Extracted from a medical imaging application (MRI processing)
- Normal Gaussian stencils smooth images, but destroy sharp edges.
- This kernel performs anistropic filtering thus preserving edges.
- We may scale the size of the stencil (radius=3,5)
  - 7<sup>3</sup>-pt or 11<sup>3</sup>-pt stencils.
  - apply to dataset of 192 x 256x256 slices
  - originally 8-bit grayscale voxels, but processed as 32-bit floats



## **3D Bilateral Filtering**

(pseudo code)

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Each point in the stencil mandates a voxel-dependent indirection, and each stencil also requires one divide.

```
for all points (xyz) in x,y,z{
  voxelSum = 0
  weightSum = 0
  srcVoxel = src[xyz]
  for all neighbors (ijk) within radius of xyz{
    neighborVoxel = src[ijk]
    neighborWeight = table2[ijk]*table1[neighborVoxel-srcVoxel]
    voxelSum +=neighborWeight*neighborVoxel
    weightSum+=neighborWeight
  }
  dstVoxel = voxelSum/weightSum
}
```

 Large radii results in extremely compute-intensive kernels with large working sets



## **Benchmark Machines**

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- Experiments only explored parallelism within an SMP
- We use a Sun X2200 M2 as a proxy for the XT5 (e.g. Jaguar)
- We use a Nehalem machine as a proxy for possible future Cray machines.
- Barcelona/Nehalem are NUMA



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# Generalized Framework for Auto-tuning Stencils

Copy and Paste auto-tuning

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Given a F95 implementation of an application:

- 1. Programmer annotates target stencil loop nests
- 2. Auto-tuning System:
  - converts FORTRAN implementation into internal representation (AST)
  - builds a test harness
  - Strategy Engine iterates on:
    - apply optimization to internal representation
    - backend generation of optimized C code
    - compile C code
    - benchmark C code
  - using best implementation, automatically produces a library for that kernel/machine combination
- 3. Programmer then updates application to call optimized library routine



## Strategy Engine: Auto-parallelization

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- The strategy engines can auto-parallelize cache blocks among hardware thread contexts.
- We use a single-program, multiple-data (SPMD) model implemented with POSIX Threads (Pthreads).
- ✤ All threads are created at the beginning of the application.

We also produce an initialization routine that exploits the first touch policy to ensure proper NUMA-aware allocation.



## Strategy Engine: Auto-tuning Optimizations

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- Strategy Engine explores a number of auto-tuning optimizations:
  - loop unrolling/register blocking
  - cache blocking
  - constant propagation / common subexpression elimination



- Future Work:
  - cache bypass (e.g. *movntpd*)
  - software prefetching
  - SIMD intrinsics
  - data structure transformations



# **Experimental Results**

NOTE: threads are ordered to exploit: multiple threads within a core (Nehalem only), then multicore, then multiple sockets (Barcelona/Nehalem)



- On the memory-bound architecture (Barcelona), auto-parallelization doesn't make a difference.
- Auto-tuning enables scalability.
- Barcelona is bandwidth-proportionally faster than the XT4.
- ✤ Nehalem is ~2.5x faster than Barcelona, and 4x faster than the XT4
- Auto-parallelization plus tuning significantly outperforms OpenMP.





- \* No changes to the framework were required (just drop in F95 code)
- As there was less reuse in the Divergence than in Laplacian, there are fewer capacity misses.
- So auto-tuning has less to improve upon
- Nehalem is ~2.5x faster than Barcelona





- \* No changes to the framework were required (just drop in F95 code)
- Gradient has moderate reuse, but a large number of output streams.
- Performance gains from auto-tuning are moderate (25-35%)
- Parallelization is only valuable in conjunction with auto-tuning





- \* No changes to the framework were required (just drop in F95 code)
- Essentially a 7x7x7 (343-pt) stencil
- Performance is much more closely tied to GHz instead of GB/s.
- Auto-parallelization yielded near perfect parallel efficiency wrt cores on Barcelona/Nehalem (Nehalem has HyperThreading)
- Auto-tuning significantly outperformed OpenMP (75% on Nehalem)





- basically the same story as radius=3
- XT4/Nehalem delivered approximately same performance as they did with radius=3
- Barcelona delivered somewhat better performance.





# Summary

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- Solution Step Forward in auto-tuning technology
- Although the framework required substantial up front work, it provides performance portability across the breadth of architectures AND stencil kernels.
- Delivers very good performance, and well in excess of OpenMP.
- Future work will examine relevant optimizations
  - e.g. cache bypass would significantly improve gradient performance.



## Summary: Machine Comparison

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- Barcelona delivers bandwidth-proportionally better performance on the memory-intensive differential operators.
- Surprisingly, Barcelona delivers ~2.5x better performance on the compute intensive bilateral filter.
- Nehalem clearly sustains dramatically better performance than either Opteron.
- Despite having a 15% faster clock, nehalem realizes a much better bilateral filter performance.





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## **Questions?**

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