CRD

BERKELEY LAB

Advanced Checkpoint Fault Tolerance Solutions for HPC

Paul Hargrove with Eric Roman and Jason Duell
checkpoint@lbl.gov

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Office of Science
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http://ftg.lbl.gov/checkpoint

Introduction

Checkpoint. Save a process's state to a file.
Restart. Reconstruct the process from a file.
BLCR. Berkeley Lab Checkpoint Restart for Linux.

Project goals. What is BLCR's approach to CR?
Why use checkpoint/restart?

System design. How does BLCR work?

Current status. What does BLCR do now?

Plans. Where is BLCR going?

Project Goals

Provide checkpoint/restart for Linux clusters running scientific workloads.

Checkpoint and restart *jobs* (shell scripts) running MPI applications.
Support a wide variety of networks.

Fit easily into production systems

Run *unmodified* application source.
Run *unmodified* binaries where possible. No special compile/link in most cases.
Run on *unpatched* kernels (as a kernel module).
Run with *unmodified* system libraries (e.g. libc).

Unrelated features (ptrace, Unix domain sockets) have low implementation priority

Why checkpoint?

We see three main scenarios: scheduling, fault tolerance and debugging.

Usage Scenarios

Batch Scheduling.

- C/R can be used to preempt and/or migrate running jobs.
- Drain queues quickly for maintenance.
- Increase system throughput by switching job mix between long jobs and wide jobs.
- Increase system utilization by allowing the scheduler to correct for bad decisions.
- Gang scheduling. Divide system time up into slots.
- Priority scheduling. Run jobs with the highest priority.

Fault Tolerance.

- Not every application can checkpoint itself.
- Periodic checkpoints can reduce lost work in case of failure (but adds cost to normal fault-free execution).
- Reactive checkpoints can respond to non-yet-fatal problems (like loss of a fan).

Debugging.

- Rollback execution to a checkpoint taken before a fault, restart with a debugger.

Other Approaches

Application-based checkpointing.

Efficient: save only needed data as step completes.
Good for fault tolerance: bad for preemptive scheduling.
Requires per-application effort by programmer.

Library-based checkpointing.

Portable across operating systems.
Transparent to application (but may require relink, etc.).
Can't (generally) restore all resources (ex: process IDs).
Can’t checkpoint shell scripts.

Hypervisor (similar arguments for software suspend).

Granularity is a full virtual machine.
Administrators have to maintain one VM per checkpoint.
Rollback. What happens to the disk state?
Debugging?
Coordination for distributed jobs is still necessary.
Scheduler integration.

Implementation

BLCR provides single node checkpoint/restart through kernel modules and a runtime library.

libcr.so: Full library: can register handlers, request checkpoints, etc.

OR *libcr_run.so*: Stub library with only a default checkpoint handler

Kernel modules: coordinates the process checkpoints, saves/restores kernel data structures, interfaces with library and command line tools.

BLCR doesn't provide built-in support for distributed runtime features

TCP sockets, bproc namespaces, etc.

Instead, BLCR provides hooks which allow apps and libraries to coordinate checkpoints and restart distributed processes through *callbacks*.

So, the MPI library must know how to checkpoint; the user application does not.

Basic Operation

Rough idea: Send the application a signal that tells it to call into BLCR.

A checkpoint request can come from the same process, or from another.
By default, user code doesn't *need* to do anything to handle it.
If desired, user code *may* register a callback to handle it.
If desired, user code *may* block requests (critical sections).

Example: Migration

gaius:~ <2>

```
pcp-x-1% ./counting
Counting demo starting with pid 3382
Count = 0
Count = 1
Count = 2
Count = 3
Count = 4
[1]    3382 killed     ./counting
pcp-x-1%
```

xterm

```
pcp-x-2% cd /home/pcp1/eroman/src/lbnl_cr/build.pcp-x-1/examples/counting
pcp-x-2% ls
context.3382  counting  counting.o  Makefile
pcp-x-2% cr_restart context.3382
Count = 5
Count = 6
Count = 7

pcp-x-2%
```

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Status

Processes, process groups and sessions

- Shell scripts (bash, tcsh, python, perl, ruby, ...).
- Multithreaded processes (pthreads with standard NPTL).
- Resources shared between processes are restored.
- Restore PID and parent PID.

Files

- Reopen files during restart: open, truncate, and seek.
- Pipes and named FIFOs.
- Files must exist in same location on filesystem.
- Memory mapped files are remapped.
- Option to save shared libraries and executable.
- Option for file path relocation.

Supported Platforms

Linux kernel 2.6 kernels

- test with kernels from kernel.org, Fedora, SuSE, and Ubuntu
- support of custom patched kernels through autoconf

Architectures

- x86, x86-64, ppc, ppc64 and ARM
- Xen dom0 and domU

MPI Implementations

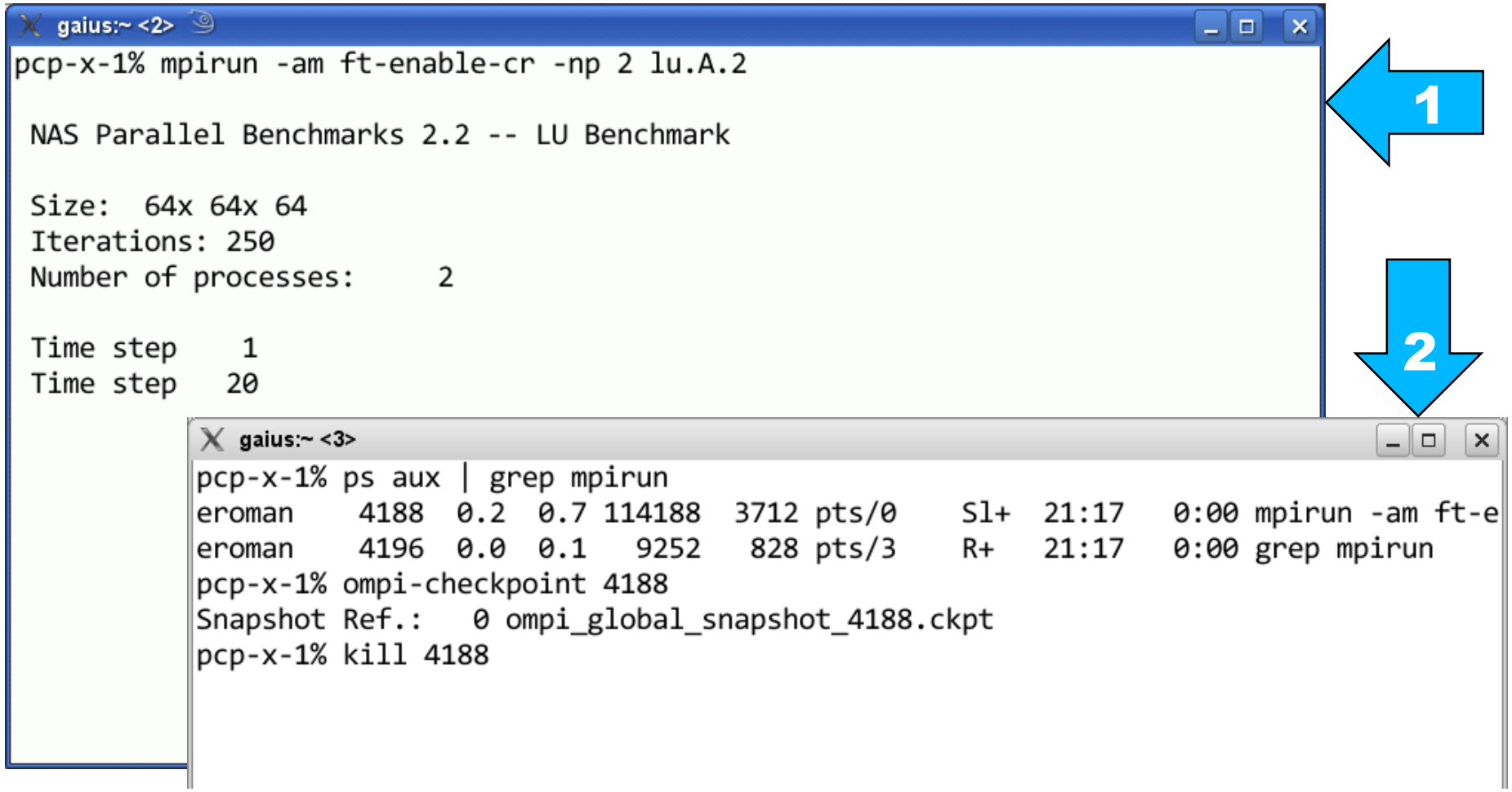
- MVAPICH2
- LAM/MPI 7.x (sockets and GM)
- MPICH-V 1.0.x with sockets
- OpenMPI
- Cray Portals

Batch Queue Systems

- Torque support available in recent snapshots.
- qhold, qrls, and periodic checkpoints tested.
- BLCR, Condor and Parrot HOWTO available.

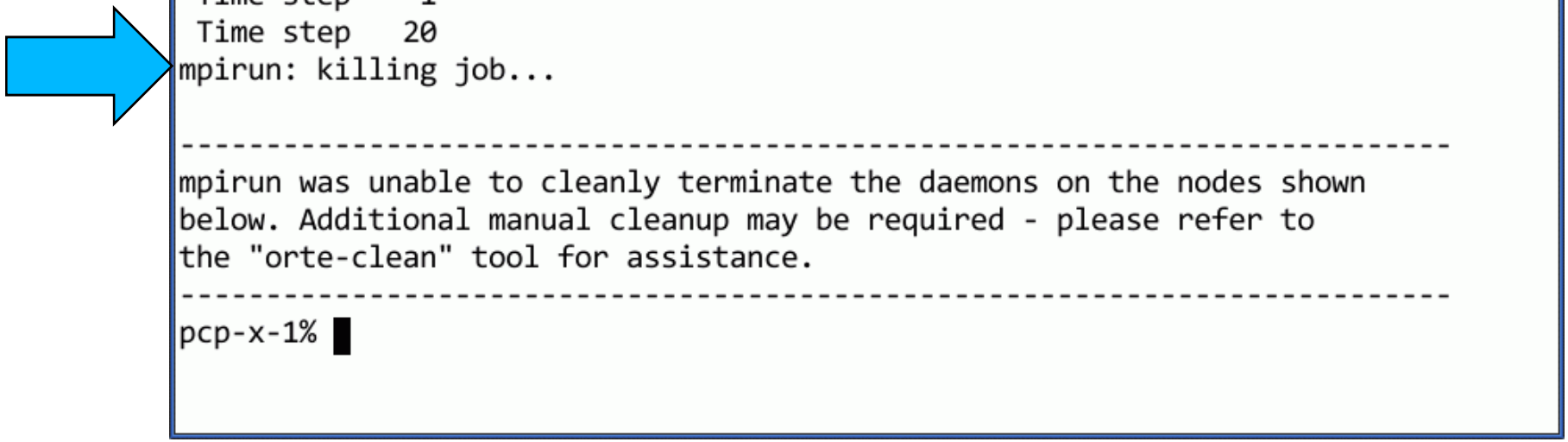
Example 2: MPI Checkpoint/Restart

Step 1 (mpirun) and Step 2 (checkpoint)



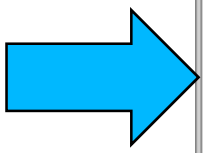
Example 2: MPI Checkpoint/Restart

The job terminates...



Example 2: MPI Checkpoint/Restart

The job restarts...



```
gaius:~ <3>
pcp-x-1% ps aux | grep mpirun
eroman    4188  0.2  0.7 114188  3712 pts/0    Sl+  21:17   0:00 mpirun -am ft-e
eroman    4196  0.0  0.1   9252   828 pts/3    R+   21:17   0:00 grep mpirun
pcp-x-1% ompi-checkpoint 4188
Snapshot Ref.:   0 ompi_global_snapshot_4188.ckpt
pcp-x-1% kill 4188
pcp-x-1% ompi-restart ompi_global_snapshot_4188.ckpt
 Time step   40
 Time step   60
 Time step   80
 Time step  100
 Time step  120
 Time step  140
 Time step  160
 Time step  180
 Time step  200
 Time step  220
```

Work In Progress

Queue system support

BLCR, Torque, and OpenMPI

Alternative handling of files

Allow checksum of file, with restart error if it has changed.

Allow saving contents of file (restore may either replace or rename)

Support files that are not open at checkpoint time, but are specified as being part of the checkpoint

Improved I/O

On-the-fly compression of context files

Direct I/O

Other

Detailed error reporting (*e.g.* What file caused ENOENT?)

Zombie processes

Conclusions

Future Work

- Interested in other queue systems (LSF, SGE, SLURM, etc.)
- More MPI implementations
- MPICH 2 support anticipated
- Vendor support (Quadrics)?
- MPI support for partial/live migration
- Ship support with distributions (ROCKS, OSCAR)

We expect BLCR to be deployed in a production batch environment before the end of the calendar year.

Torque support will be available soon.

You should be able to install BLCR on your system and checkpoint your MPI applications with it.

We would like you to download BLCR and try it!

For More Information

http://ftg.lbl.gov/checkpoint

Papers (available from website):

- "*Design and Implementation of BLCR*": high-level system design, including description of user API
- "*Requirements for Linux Checkpoint/Restart*": exhaustive list of Unix features we will support (or not).
- "*A Survey of Checkpoint/Restart Implementations*": focusing on open source versions that run on Linux
- "*The LAM/MPI Checkpoint/Restart Framework: System-Initiated Checkpointing*": implementation with LAM/MPI

CIFTS

Coordinated Infrastructure for Fault Tolerant Systems

Parent project. Building a notification infrastructure for BLCR.

http://www.mcs.anl.gov/research/cifts/